

# ENERGY AWARE QUALITY OF SERVICE ROUTING PROTOCOL FOR LARGE SCALE MANET

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**ABSTRACT:** *One of the important and challenging problems in the design of ad hoc networks is the development of an efficient routing protocol that can provide high-quality communications among mobile hosts for that proposing new protocol to evaluate the node lifetime and the link lifetime utilizing the dynamic nature, such as the energy drain rate and the relative mobility estimation rate of nodes. Integrating these two performance metrics by using the proposed route lifetime-prediction algorithm select the least dynamic route with the longest lifetime for persistent data forwarding and based on quadrant. Our proposed route Lifetime-prediction protocol in a exploring dynamic nature routing for large scale network (LEDNR) protocol environment based on Adhoc on demand distance vector routing (AODV)*

**Keywords:** *Mobile ad hoc networks (MANETs), node lifetime, route discovery, Proactive/reactive routing protocol, performance analysis.*

## 1. INTRODUCTION

A mobile ad hoc network (MANET) consists of many mobile nodes that can communicate with each other directly or through intermediate nodes. Often, hosts in a MANET operate with batteries and can roam freely, and thus, a host may exhaust its power or move away, giving no notice to its neighboring nodes, causing changes in network topology. A key characteristic of these scenarios is the dynamic behaviour of the involved communication partners.

Communication protocols will have to deal with a frequently changing network topology. However, many applications require stable connections to guarantee a certain degree of QoS. In

access networks, access point handovers may disrupt the data transfer. In addition, service contexts may need to be transferred to the new access points, introducing additional overhead and delays to the connection. In ad hoc networks, mobile services enable peer-to-peer connections for voice or data traffic. Using stable links is crucial for establishing stable paths between connection peers. Rerouting is especially costly in these networks without infrastructure, since it usually results in (at least partly) flooding the network. The stability of a link is given by its probability to persist for a certain time span, In MANETs, a route consists of multiple links in series, and thus, its lifetime depends on the lifetime of each node, as well as the wireless links between adjacent nodes. The main contribution of this paper is that we combine node lifetime and LLT in route lifetime-prediction algorithm, which explores the dynamic nature of mobile nodes the energy drain rate of nodes and the relative mobility estimation rate at which adjacent nodes move apart) in a route-discovery period that predicts the lifetime of routes discovered, and then, we select the longest lifetime route for persistent data forwarding when making a route decision. The proposed route lifetime-prediction algorithm. Implements our proposed algorithm in an exploring dynamic nature routing (LEDNR) protocol large scale environment based on quadrant based dynamic source routing.

## **2.PROBLEM DEFINITION**

1. Any one of the nodes in the route dies because of limited battery energy
2. Any one of the connections is broken because the corresponding two adjacent nodes move out of each other's communication range. Thus the lifetime of route P is expressed as the minimum value of the lifetime of both nodes and connections involved in route.
3. By considering the energy state of nodes, such as residual energy and energy drain rate, the node lifetime routing algorithms often select a path consisting of nodes that may survive for the longest time among Multiple paths
4. The energy drain rate of a node is affected not only by its own but by its neighboring data flows as well.
5. Aimed to preserve network connectivity by choosing a route according to the remaining battery.
6. Life of nodes along the route proposed selecting a path with minimum total transmission power when there exist some possible paths, and all nodes through these paths have sufficient residual battery power
7. Selecting a path that has the largest packet transmission capacity (the residual energy divided by the expected energy spent in reliably forwarding a packet) at a "critical" node among multiple

paths. The critical node is the node that has the smallest packet transmission capacity in a path. Each node attempts to estimate its battery lifetime based on its residual energy and its past activity all of them ignored the mobility of mobile hosts, and thus, it seems that they are more suitable for static networks.

### 3. THE PRINCIPLE OF LEDNR METHODOLOGY:

The proposed algorithm consists of the following three phases: **Route discovery, Data forwarding, and Route maintenance.**

There are seven main differences between the LEDNR and the AODV. **First**, in the LEDNR protocol, every node saves the received signal strength and the received time of the RREQ packet in its local memory and adds this information into the RREP packet header in a piggyback manner when it receives the RREP for the corresponding RREQ packet to meet the requirement of the connection lifetime-prediction algorithm. **Second**, node agents need to update their predicted node lifetime.

during every period. **Third** the node-lifetime information in the RREP packet is updated when the RREP packet is returned from a destination node to the source node. **Fourth** The location information of the source and destination nodes is piggy-backed in the route request (RREQ) packet and then broadcasted. Upon receiving the RREQ, intermediate nodes will compare using a simple mathematical comparison based on the coordinates of source, destination and the current node that directs the packet towards the destination Once the decision to broadcast has been made **Fifth** the intermediate node will insert its location by replacing the source node coordinates and append its address and sequence number at the end of the RREQ packet.

It will then broadcast the packet. **Sixth** The process will repeat at each intermediate node until it reaches the destination. The replacement of the source node location information with the intermediate node coordinates will make the packet more directed towards the destination since the comparison now is based on the previous node. **Finally** Upon receiving the RREQ, destination node will send a route reply message (RREP) back to source via the path taken to reach the destination that was appended in the RREQ as it traverses across the network. There is no need for the route discovery to the source node. Figure1 shows the format of the RREQ packet in Q-DIR where the source and destination nodes location information are inserted are highlighted.

**Node Lifetime Prediction Algorithm:**

A link is composed of the two nodes in a connection and the connection itself, and the LLT includes both the node lifetime and the connection lifetime. A link  $L_i$  consists of a connection  $C_i$  and two nodes  $(N_{i-1}, N_i)$ .

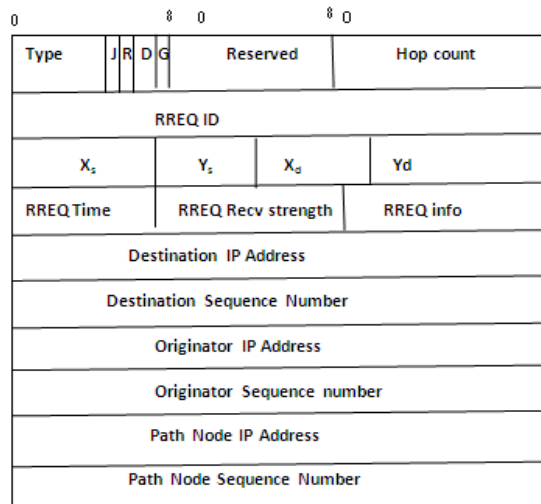


Fig1:RREQ control message format of LEDNR

where  $C_i$  represents the connection between nodes  $N_{i-1}$  and  $N_i$  and it is maintained until the adjacent nodes  $(N_{i-1}, N_i)$  move out of each other’s communication range under the assumption of no energy problem in both nodes  $N_{i-1}$  and  $N_i$ . We introduce connection lifetime  $TC_i$  to represent the estimated lifetime of the connection  $C_i$ , and it only depends on their relative mobility and distance of nodes  $N_{i-1}$  and  $N_i$  at a given time. The term  $TN_i$  denotes the estimated battery lifetime of node  $N_i$ . Then, the lifetime of the link  $L_i$  is expressed as the minimum value of  $(TC_i, TN_{i-1}, TN_i)$ ,

$$TL_i = \Psi \min(TC_i, TN_{i-1}, TN_i)$$

The lifetime of route P is expressed as the minimum value of the lifetime of both nodes and connections involved in route P. Assume that  $\Omega$  represents the set of all nodes in route P and that  $\Psi$  is the set of all the connections in route P. Thus, the lifetime  $Tp$  of route P can be expressed as

$$Tp = \min(TN_i, TC_i)$$

We use an **Exponentially Weighted Moving Average Method** to estimate the energy drain rate  $ev$  evaluate the node lifetime that is based on its current residual energy and its past activity.

Where

**Eu/v** = energy spent at node u due to node v

**ETack** = energy spent for transmission of an ACK packet

**ETpck** = energy spent for transmission of a data packet

**ERack** = energy spent for reception of an ACK packet

**ERpck** = energy spent for reception of a data packet, 1 if p is true, 0 otherwise.

The total energy spent at a node is the sum of the energy spent at it due to all the nodes in the reception and interference area of this node. Thus, in this model, the transmission and reception costs are included if the node belongs to a flow, and reception costs are included if it is near the flow. Through the cost value we can predict the optimal path in dynamic nature and estimate the reliable path for the large scale network. Constant value with a range of [0, 1] the current condition of node  $i$  well, we grant a higher priority. The path lifetime (PLT), RREQ time, and RREQ signal strength, is added to the common header of an RREP packet.

The PLT represents the predicted lifetime of the source route in this packet header and can be updated when RREP packets are forwarded from the destination node to the source node in the route-discovery phase.

#### **Connection LifeTime Prediction Algorithm:**

Handle the connections that are in an unstable state and may only last for a short period particularly for ignoring the stable one for simplicity. The reasons are given as follows: First, we are only concerned with the minimum node lifetime or the connection life time in a route. Since two nodes of a stable connection are within the communication range of each other, the connection lifetime may last longer, and they are not a bottleneck from the route to which they belong. it is easier to model the mobility of nodes in a short period during which unstable connections last. We can assume reasonably and simply that the nodes move at a constant speed toward the same direction in such a short period measure the distance between nodes  $N_i$  and  $N_{i-1}$ . we use Global-Positioning-System-based location information to measure the received signal strength. Assuming that senders transmit packets with the same power level a receiver can measure the received signal power strength when receiving a packet and then calculates the distance by directly applying the radio propagation model. If the received signal power strength is lower than a threshold value, we regard this link as an unstable state and then calculate the connection time.

Our proposed method requires only two sample packets, and we implement piggyback information on route-request (RREQ) and route-reply (RREP) packets during a route-discovery procedure with no other control message overhead, and thus, it does not increase time complexity. LEDNR searches for the route whenever there is a need. In addition to the above three, DATA packet is also there and it is unicast/Broadcast. RERR messages will be assigned unconditional overhearing. The reason is that the link failure should be informed to all the nodes, so that the nodes will not use it for the next time until it gets ready RREQ is a broadcast and here comes the probability values. Based on the probability of overhearing and rebroadcast, it is set. Conditional overhearing is done for the RREP. RREQ is set to overhear under a condition. This avoids redundancy in broadcasting.

**ALGORITHM:****Predict** Its lifetime**If** its lifetime > Min-lifetime**Replace** Min-lifetime with its lifetime**If** Sequence Number exists**Compare** Min-lifetime of current RREQ with Minlifetimeof existing one.**If** new Min-lifetime <= old Min-lifetime**Discard** new RREQ**If** new Min-Lifetime >old Min-lifetime**Replace** old Min-Lifetime with new Minlifetime**4. EXPERIMENTAL SETUP**

Our implementation is based on network simulator ns2.34. The mobility of nodes follows a random way-point model. The source–destination connection patterns are generated using cbrgen.tcl in NS-2. The initial energy is the 1000 joules and simulation time is set to 100 s In the simulation, we consider a total of 1000 nodes initially randomly distributed over a square network

of size  $1000m \times 1000m$ . Every node moves at a speed  $V$  and transmits at uniform power of coverage of radius  $R$  under certain traffic load. Three different transmission ranges  $R \{150, 200, 250\}m$  are covered, all within the coverage. Four different speeds  $V \{5, 10, 15, 20\}m/s$  are simulated, from lower mobility to higher mobility scenarios. Traffic, supplied from a CBR source with fixed packet size of 1000 Bytes, is randomly generated with uniformly distributed sources and destinations. Different number of traffic flows  $F \{5, 10, 15, 20\}$  are simulated, covering low and moderate flow configuration.

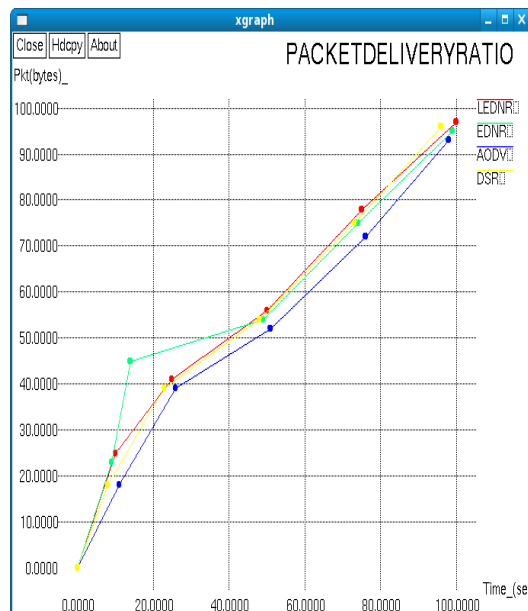
**PACKET DELIVERY RATIO**

Packet Delivery Ratio (PDR) is calculated by dividing the number of packets received by the destination through the number of packets originated by the source.

$$\frac{\text{TOTAL NUMBER OF PACKETS SUCCESSFULLY DELIVERED}}{\text{TOTAL NUMBER OF PACKETS SENT}} \times 100\%$$

Fig. 1 This result also shows us that LEDNR has a increase the packet delivery fraction than the other protocols

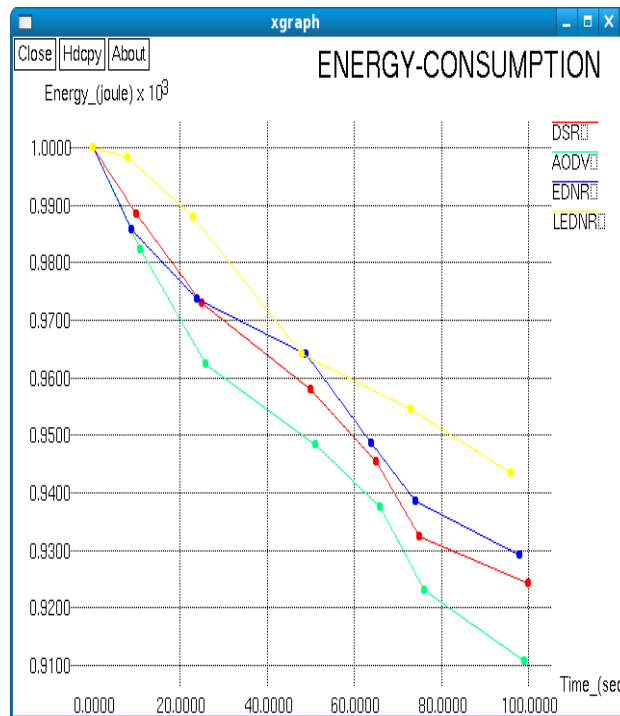
Fig 1: Packet delivery Ratio



**ENERGY CONSUMPTION:**

This measures the energy expended per delivered data packet. It is expressed as

$$\frac{\sum \text{ENERGY EXPENDED BY EACH NODE}}{\text{TOTAL NUMBER OF PACKETS DELIVERED}}$$



**Fig2** : Comparison of AODV ,DSR and EDNR protocols with LEDNR

Fig. 2 This result also shows us that LEDNR has a lesser energy consumption than the other protocols

**END-END LATENCY:**

**End-End latency** measures the average time it takes to route a data packet from the source node to the destination. It is expressed as

$$\frac{\sum \text{INDIVIDUAL DATA PACKET LATENCY}}{\text{TOTAL NUMBER OF PACKETS DELIVERED}}$$



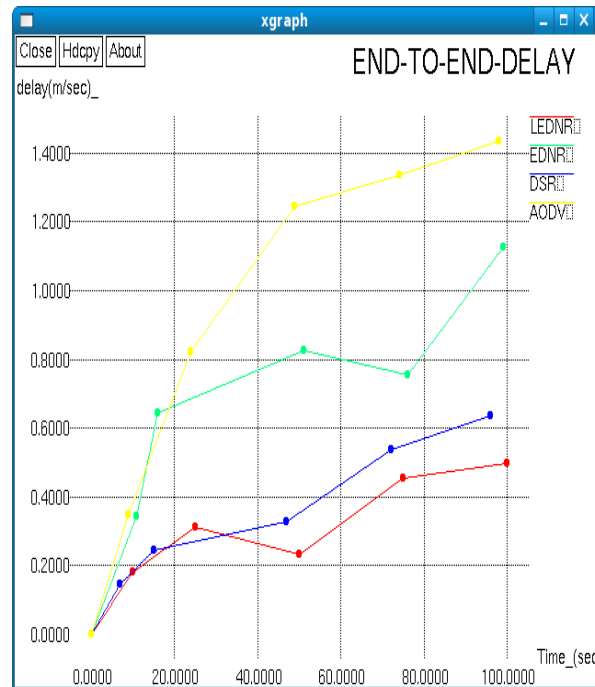


Fig. 3 This result also shows us that LEDNR has a lesser Delay than the other protocols

The advantage of the LEDNR protocol in terms of the number of routing failures delay, energy consumption. To adapt to dynamically varying network topology environments, the EDNR, AODV and DSR protocols do their best to find a more stable route.

## 5. CONCLUSION

In MANETs, a link is formed by two adjacent mobile nodes, which have limited battery energy and can roam freely, and the link is said to be broken if any of the nodes dies because they run out of energy or they move out of each other's communication range. In this paper, we have considered both the node lifetime and the LLT to predict the route lifetime and have proposed a new algorithm that explores the dynamic nature of mobile nodes, such as the energy drain rate and the relative motion estimation rate of nodes, to evaluate the node lifetime and the LLT.

Combining these two metrics by using our proposed route lifetime-prediction algorithm we can select the least dynamic route with the longest lifetime for persistent data forwarding. Finally, we have evaluated the compare performance of the proposed LEDNR protocol based on the DSR and AODV.

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